

For suppose that $g(p) = i$. Then $f(p, i) = o$, and for any $j < i, f(p, j) \neq o$. Since A represents f in Q , we have

$$\vdash_Q A(p, i, o), \text{ and (if } i > o), \quad (i)$$

$$\vdash_Q A(p, o, o), \quad (o)$$

$$\vdots \quad (i-1)$$

$$\vdash_Q A(p, i-1, o). \quad (i-1)$$

(o), ..., (i-1), and 14.11 entail that

$$\vdash_Q \forall w (w < i \rightarrow \neg A(p, w, o)), \quad (i+1)$$

which, together with (i), entails that $\vdash_Q B(p, i)$.

We must show that $\vdash_Q \forall x_{n+1} (B(p, x_{n+1}) \rightarrow x_{n+1} = i)$. Assume $B(p, x_{n+1})$, i.e., $A(p, x_{n+1}, o) \& \forall w (w < x_{n+1} \rightarrow \neg A(p, w, o))$. From (i) and

$$\forall w (w < x_{n+1} \rightarrow \neg A(p, w, o)), \text{ we have } \neg i < x_{n+1}.$$

From $A(p, x_{n+1}, o)$ and (i+1), we have $\neg x_{n+1} < i$. Thus by 14.13 we have $x_{n+1} = i$. So $\vdash_Q \forall x_{n+1} (B(p, x_{n+1}) \rightarrow x_{n+1} = i)$.

Exercises

14.1 Verify the following assertion: all recursive functions are representable in the theory ('R') whose language is L and whose theorems are the consequences in L of the following infinitely many sentences:

$$i \neq j \text{ for all } i, j \text{ such that } i \neq j;$$

$$i+j = k \text{ for all } i, j, k \text{ such that } i+j = k;$$

$$i \cdot j = k \text{ for all } i, j, k \text{ such that } i \cdot j = k;$$

$$\forall x (x < i \rightarrow x = o \vee \dots \vee x = i-1) \text{ for all } i;$$

$$\text{and } \forall x (x < i \vee x = i \vee i < x), \text{ for all } i.$$

14.2 Show that none of the following sentences are theorems of Q :

(a) $\forall x x \neq x'$,

(b) $\forall x \forall y \forall z x+(y+z) = (x+y)+z$,

(c) $\forall x \forall y x+y = y+x$,

(d) $\forall x o+x = x$,

(e) $\forall x x < x'$,

(f) $\forall x \forall y \neg (x < y \& y < x)$,

(g) $\forall x \forall y \forall z x \cdot (y \cdot z) = (x \cdot y) \cdot z$,

(h) $\forall x \forall y x \cdot y = y \cdot x$,

(i) $\forall x o \cdot x = o$,

(j) $\forall x \forall y \forall z x \cdot (y+z) = x \cdot y + x \cdot z$.

Hint: Let a and b be two objects that are not natural numbers, and consider the following successor, addition, and multiplication tables:

x	x'	$\overline{f}+$	j	a	b	$\overline{f} \cdot$	o	$j \neq o$	a	b
i	i'	i	$i+j$	b	a	o	o	o	a	b
a	a	a	a	b	a	$i \neq o$	o	$i \cdot j$	a	b
b	b	b	b	b	a	a	o	b	b	b
						b	o	a	a	a